

BUUCTF reverse: [GXYCTF2019]luck_guy,findit,简单注册器题解

原创

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本文链接: https://blog.csdn.net/weixin_50549897/article/details/113523762

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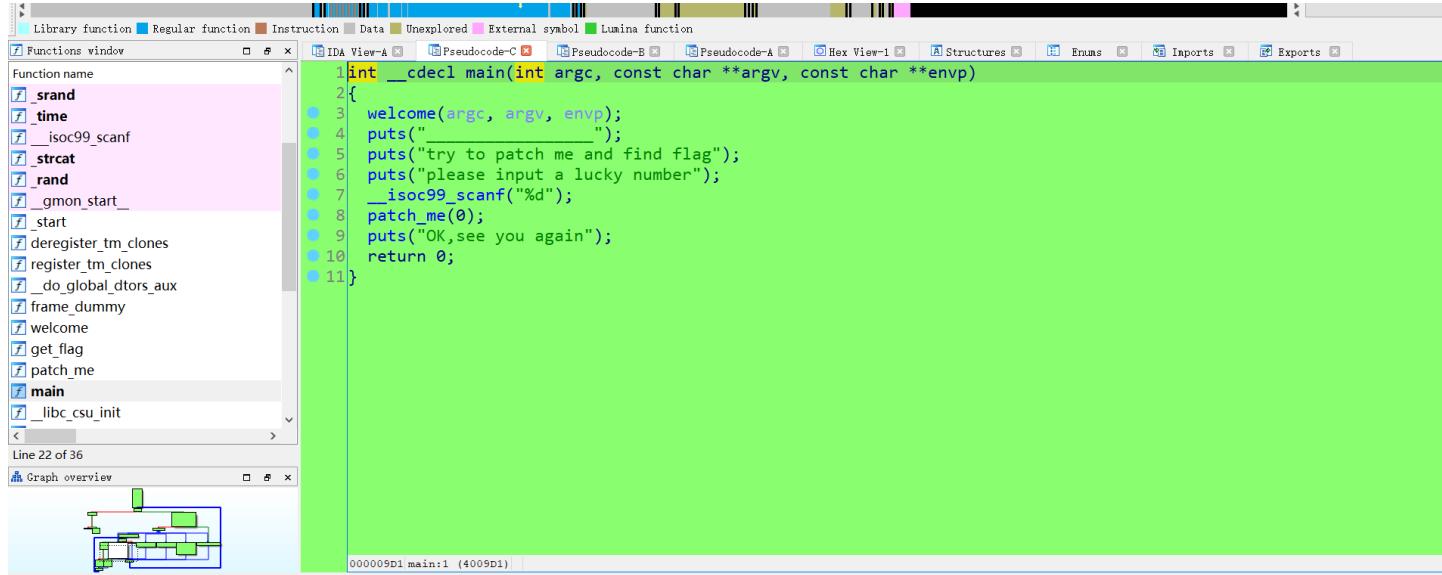
[GXYCTF2019]luck_guy, 简单注册器, 题解

由于作者水平有限, 有不当之处或有更好方法请在评论区指出

题目地址<https://buuoj.cn/challenges>

luck_guy

拖进IDA, 查看main函数



```
int __cdecl main(int argc, const char **argv, const char **envp)
{
    welcome(argc, argv, envp);
    puts("____");
    puts("try to patch me and find flag");
    puts("please input a lucky number");
    __isoc99_scanf("%d");
    patch_me(0);
    puts("OK,see you again");
    return 0;
}
```

The screenshot shows the IDA Pro interface with the main function decompiled. The code is as follows:

```
int __cdecl main(int argc, const char **argv, const char **envp)
{
    welcome(argc, argv, envp);
    puts("____");
    puts("try to patch me and find flag");
    puts("please input a lucky number");
    __isoc99_scanf("%d");
    patch_me(0);
    puts("OK,see you again");
    return 0;
}
```

The 'Functions window' on the left lists several functions, with 'main' currently selected. The 'Graph overview' window at the bottom shows the control flow graph for the main function.

找到patch_me(0), 里面可以看到, 如果输入的是偶数就进入get_flag()函数。本题关键应该就在里面。打开后

```

unsigned __int64 get_flag()
{
    unsigned int v0; // eax
    int i; // [rsp+4h] [rbp-3Ch]
    int j; // [rsp+8h] [rbp-38h]
    __int64 s; // [rsp+10h] [rbp-30h] BYREF
    char v5; // [rsp+18h] [rbp-28h]
    unsigned __int64 v6; // [rsp+38h] [rbp-8h]

    v6 = __readfsqword(0x28u);
    v0 = time(0LL);
    srand(v0);
    for ( i = 0; i <= 4; ++i )
    {
        switch ( rand() % 200 ) // 产生1-199之间的随机数
        {
            case 1:
                puts("OK, it's flag:");
                memset(&s, 0, 0x28uLL);
                strcat((char *)&s, f1);
                strcat((char *)&s, &f2); // f2是空的
                printf("%s", (const char *)&s);
                break;
            case 2:
                printf("Solar not like you");
                break;
            case 3:
                printf("Solar want a girlfriend");
                break;
            case 4:
                s = 0x7F666F6067756369LL;
                v5 = 0;
                strcat(&f2, (const char *)&s); // 将f2与s拼接
                break;
            case 5:
                for ( j = 0; j <= 7; ++j )
                {
                    if ( j % 2 == 1 )
                        *(&f2 + j) -= 2;
                    else
                        --*(&f2 + j);
                }
                break;
            default:
                puts("emmm,you can't find flag 23333");
                break;
        }
    }
    return __readfsqword(0x28u) ^ v6;
}

```

分析伪代码。从rand() % 200可以看出来，随机生成1-199之间的数字。表示这些书随机组合，flag就在其中。原来题目luck_guy是这个意思。显然，flag由f1和f2拼接而成。我们可以看到f1为GXY{do_not_

<code>__strcat</code>	<code>.data:0000000000601077</code>	<code>db 0</code>
<code>__rand</code>	<code>.data:0000000000601078</code>	<code>public f1</code>
<code>__gmon_start_</code>	<code>.data:0000000000601078 ; char f1[]</code>	
<code>__start</code>	<code>.data:0000000000601078 f1</code>	<code>db 'GXY{do_not_',0</code>
	<code>.data:0000000000601078 .data</code>	<code>; DATA XREF: get_flag+9E1o ends</code>

选择关键是找出f2是多少。经判断，switch里的顺序是case4,case5,case1

写python脚本如下

```
flag = 'GXY{do_not_'
f2 = [0x7F, 0x66, 0x6F, 0x60, 0x67, 0x75, 0x63, 0x69]
s = ''
for i in range(8):
    if i % 2 == 1:
        c = chr(int(f2[i]) - 2)
    else:
        c = chr(int(f2[i]) - 1)
    flag += c
print(flag)
```

发现输出为GXY{do_not_~dn^fsbg，显然不是想要的flag

查资料想到IDA反编译成C语言代码，`s = 0x7F666F6067756369LL`数据储存应该是小端储存。关于大小端储存详见这篇文章所以应该反过来。

```
flag = 'GXY{do_not_'
f2 = [0x7F, 0x66, 0x6F, 0x60, 0x67, 0x75, 0x63, 0x69][::-1]
s = ''
for i in range(8):
    if i % 2 == 1:
        c = chr(int(f2[i]) - 2)
    else:
        c = chr(int(f2[i]) - 1)
    flag += c
print(flag)
```

```
GXY{do_not_hate_me}
[Finished in 1.2s]
```

简单注册器

用 android killer 打开，使用工具查看java源码（以下截取了一段）

```

if (j == 1)
{
    paramAnonymousView = "dd2940c04462b4dd7c450528835cca15".toCharArray();
    paramAnonymousView[2] = ((char)(paramAnonymousView[2] + paramAnonymousView[3] - 50));
    paramAnonymousView[4] = ((char)(paramAnonymousView[2] + paramAnonymousView[5] - 48));
    paramAnonymousView[30] = ((char)(paramAnonymousView[31] + paramAnonymousView[9] - 48));
    paramAnonymousView[14] = ((char)(paramAnonymousView[27] + paramAnonymousView[28] - 97));
    j = 0;
    for (;;)
    {
        if (j >= 16)
        {
            paramAnonymousView = String.valueOf(paramAnonymousView);
            localTextView.setText("flag{" + paramAnonymousView + "}");
            return;
        }
        int i = paramAnonymousView[(31 - j)];
        paramAnonymousView[(31 - j)] = paramAnonymousView[j];
        paramAnonymousView[j] = i;
        j += 1;
    }
}
localTextView.setText("输入注册码错误");
}
});
}

```

由此可见flag就是dd2940c04462b4dd7c450528835cca15经过了一系列加密后的结果，加密操作要进行16次写脚本如下

```

s=['d','d','2','9','4','0','c','0','4','4','6','2','b','4','d','d','7','c','4','5','0','5','2','8','8','3','5','c','c','a','1','5']
s[2]=chr(ord(s[2])+ord(s[3])-50)
s[4]=chr(ord(s[2])+ord(s[5])-48)
s[30]=chr(ord(s[31])+ord(s[9])-48)
s[14]=chr(ord(s[27])+ord(s[28])-97)

for j in range(16):
    i=s[31-j]
    s[31-j]=s[j]
    s[j]=i

for j in s:
    print(j,end="")

```

输出结果

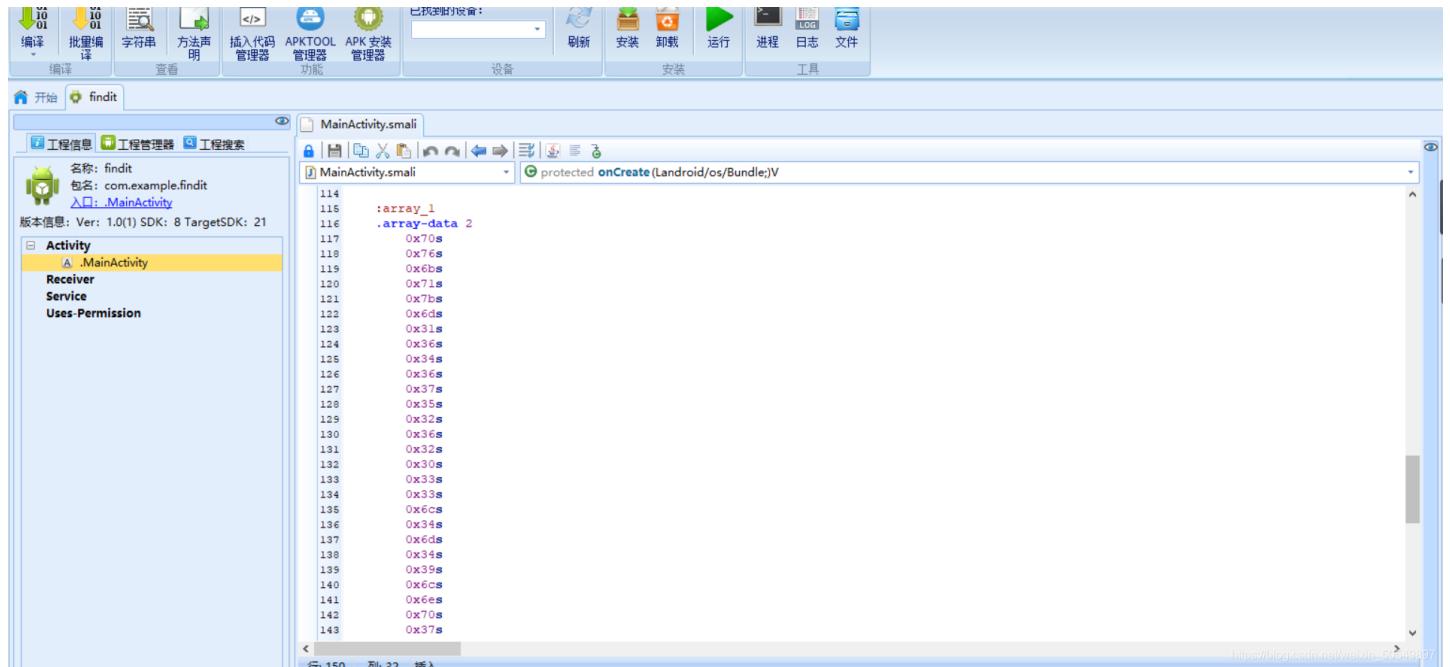
59acc538825054c7de4b26440c0999dd [Finished in 0.8s]

https://blog.csdn.net/weixin_50549897

Get the flag! flag{59acc538825054c7de4b26440c0999dd}

findit

用android killer打开,看到一组16进制数字



转换为字符后为pvkq{m164675262033l4m49lnp7p9mnk28k75}

而前面pvkq可能是flag经过某种变换加密后的结果

疑似是flag进行了某种移位，猜测是凯撒密码

放在在线解密平台上一个个密钥试看看可以找到flag

凯撒密码

Caesar Cipher

pvkq{m164675262033l4m49lnp7p9mnk28k75}

10 移除标点 (Remove Punctuation)

加密 **解密**

flag{c164675262033b4c49bdf7f9cda28a75}

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当然也可以写脚本批量输出

解密python脚本如下，把每个偏移量都遍历一遍

```
#凯撒密码破解
s='pvkq{m16467526203314m49lnp7p9mnk28k75}'
str='abcdefghijklmnopqrstuvwxyz'

for i in range(1,27):#i是密钥
    print("Key=%d"%i,end=" ")
    for word in s:
        if word in str:#如果word是字母就进行操作
            n=str.find(word)
            word=str[n-i]

        print(word,end="")
    print('\n')
```

输出的结果如下

```
Key=1 oujp{l1164675262033k4149kmo7o91mj28j75}

Key=2 ntio{k164675262033j4k49jln7n9kli28i75}

Key=3 mshn{j164675262033i4j49ikm7m9jkh28h75}

Key=4 lrgm{i1164675262033h4i49hjl7l9ijg28g75}

Key=5 kqfl{h164675262033g4h49gik7k9hif28f75}

Key=6 jpek{g164675262033f4g49fhj7j9ghe28e75}

Key=7 iodj{f164675262033e4f49egi7i9fgd28d75}

Key=8 hncl{e164675262033d4e49dfh7h9efc28c75}

Key=9 gmbh{d164675262033c4d49ceg7g9deb28b75}

Key=10 flag{c164675262033b4c49bdf7f9cda28a75}

Key=11 ekzf{b164675262033a4b49ace7e9bcz28z75}

Key=12 djye{a164675262033z4a49zbd7d9aby28y75}

Key=13 cixd{z164675262033y4z49yac7c9zax28x75}

Key=14 bhwc{y164675262033x4y49xzb7b9yzw28w75}

Key=15 agvb{x164675262033w4x49wya7a9xyv28v75}

Key=16 zfua{w164675262033v4w49vxz7z9wxu28u75}

Key=17 yetz{v164675262033u4v49uwy7y9vwt28t75}

Key=18 xdsy{u164675262033t4u49tvx7x9uvs28s75}

Key=19 wcrx{t164675262033s4t49suw7w9tur28r75}

Key=20 vbqw{s164675262033r4s49rtv7v9stq28q75}

Key=21 uapv{r164675262033q4r49qsu7u9rsp28p75}

Key=22 tzou{q164675262033p4q49prt7t9qro28o75}

Key=23 synt{p164675262033o4p49oqs7s9pqn28n75}

Key=24 rxms{o164675262033n4o49npr7r9opm28m75}

Key=25 qwlr{n164675262033m4n49moq7q9nol28l75}

Key=26 pvkq{m164675262033l4m49lnp7p9mnk28k75}

[Finished in 0.6s]
```

找到flag为flag{c164675262033b4c49bdf7f9cda28a75}密钥为10.

本人其它文章链接

[BUUCTF reverse: \[GXYCTF2019\]luck_guy,findit,简单注册器题解](#)

[封神台靶场尤里的复仇I第一第二第五第六第七章解题思路\(持续更新\)](#)

[ctfhub:网鼎杯第一场2018 reverse-beijing题解](#)

[逆向工程入门：IDAwindows本地动态调试，linux远程动态调试及虚拟机配置](#)

[逆向迷宫题总结（持续更新）2020华南师大CTF新生赛maze，攻防世界新手区：NJUST CTF 2017，BUUCTF：不一样的flag](#)