

# 2018年强网杯初赛 逆向题目 hide writeup (超详细)

原创

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订阅专栏

这道题有壳, strings搜索是upx3.91。但是不能用upx -d。只好手动脱壳了。

难点在于本题目有反调试, 有些函数不能步过。

首先研究壳。

```
LOAD:000000000044EFB0          public start
LOAD:000000000044EFB0 start      proc near
LOAD:000000000044EFB0           call loc_44F230 ; DATA XREF: LOAD:0000000000400018↑o
LOAD:000000000044EFB0           push rbp
LOAD:000000000044EFB5           push rbx
LOAD:000000000044EFB6           push rcx
LOAD:000000000044EFB7           push rdx
LOAD:000000000044EFB8           push rsi, rdi
LOAD:000000000044EFB9           push rsi
LOAD:000000000044EFBC          mov rsi, rdi
LOAD:000000000044EFBD          mov rdi, rdx
LOAD:000000000044EFC0          xor ebx, ebx
LOAD:000000000044EFC3          xor ecx, ecx
LOAD:000000000044EFC5          or rbp, 0xFFFFFFFFFFFFFFFh
LOAD:000000000044EFC7          call sub_44F020
LOAD:000000000044EFCB          add ebx, ebx
LOAD:000000000044EFD0           jz short loc_44EFD6
LOAD:000000000044EFD2           rep retn
LOAD:000000000044EFD4
LOAD:000000000044EFD6           ; -----
LOAD:000000000044EFD6 loc_44EFD6: ; CODE XREF: start+22↑j
LOAD:000000000044EFD6           mov ebx, [rsi]
LOAD:000000000044EFD8           sub rsi, 0xFFFFFFFFFFFFFFFCh
LOAD:000000000044EFDC          adc ebx, ebx
LOAD:000000000044EFD8           mov dl, [rsi]
LOAD:000000000044EFE0           rep retn
LOAD:000000000044EFE0 start    endp ; sp-analysis failed
I NAD:000000000044FFFFA
```

第一句call必须步入。

```

LOAD:00000000004F230 loc_44F230:          ; CODE XREF: start↑p
LOAD:00000000004F230    pop    rbp
LOAD:00000000004F231    lea    rax, [rbp-9]
LOAD:00000000004F235    mov    r15d, [rax]
LOAD:00000000004F238    mov    edx, 0C8h
LOAD:00000000004F23D    sub    rax, r15
LOAD:00000000004F240    sub    r15d, edx
LOAD:00000000004F243    lea    rcx, [rax+rdx]
LOAD:00000000004F247    call   sub_44F1C5
LOAD:00000000004F24C    jb    short loc_44F254
LOAD:00000000004F24C    ; -----
LOAD:00000000004F24E    dw    0
LOAD:00000000004F250    db    6Ch, 5, 2 dup(0)
LOAD:00000000004F254    ; -----
LOAD:00000000004F254 loc_44F254:          ; CODE XREF: LOAD:00000000004F24C↑j
LOAD:00000000004F254    or     [rcx+19h], cl
LOAD:00000000004F257    add    dh, bh
LOAD:00000000004F257    ; -----
LOAD:00000000004F259    db    2 dup(0FFh), 0E6h, 0E8h, 5Fh, 0, 2Fh
LOAD:00000000004F260 aProcSelfExe      db    'proc/self/exe', 0
... .procselfexe

```

Sub\_44f1c5必须步入。

Sub\_44F1C5中，

```

LOAD:00000000004F1E5    pop    r10      ; flags
LOAD:00000000004F1E7    sub    r8d, r8d  ; fd
LOAD:00000000004F1EA    push   9
LOAD:00000000004F1EC    pop    rax
LOAD:00000000004F1ED    syscall           ; LINUX - sys_mmap
LOAD:00000000004F1EF    cmp    edi, eax
LOAD:00000000004F1F1    jnz    loc_44F0EB
LOAD:00000000004F1F7    mov    esi, offset dword_400000
LOAD:00000000004F1FC    mov    edx, edi
LOAD:00000000004F1FE    sub    edx, esi
LOAD:00000000004F200    jz    short loc_44F217
LOAD:00000000004F202    add    ebp, edx  ; 增加偏移的地址差
LOAD:00000000004F204    add    [rsp+28h+var_20], edx
LOAD:00000000004F208    add    [rsp+28h+var_18], edx
LOAD:00000000004F20C    ; -----
LOAD:00000000004F20C loc_44F20C:          ; CODE XREF: LOAD:00000000004F1A5↑j
LOAD:00000000004F20C    mov    ecx, ebx
LOAD:00000000004F20E    sub    ecx, esi
LOAD:00000000004F210    shr    ecx, 3
LOAD:00000000004F213    cld
LOAD:00000000004F214    rep    mousq
LOAD:00000000004F217 loc_44F217:          ; CODE XREF: sub_44F1C5+3B↑j
LOAD:00000000004F217    xchg   eax, edi
LOAD:00000000004F218    mov    rsi, rbx
LOAD:00000000004F218    push   rax
LOAD:00000000004F21C    xchg   eax, edx
LOAD:00000000004F21D    lodsd
LOAD:00000000004F21E    push   rax
LOAD:00000000004F21F    ; -----
LOAD:00000000004F21F loc_44F21F:          ; CODE XREF: LOAD:00000000004F1A8↑j
LOAD:00000000004F21F    mov    rcx, rsp
LOAD:00000000004F222    lodsd
LOAD:00000000004F223    xchg   eax, edi
LOAD:00000000004F224    lodsd
LOAD:00000000004F225    movzx  r8d, al
LOAD:00000000004F229    xchg   rdi, rsi
LOAD:00000000004F22C    call   rbp
LOAD:00000000004F22E    pop    rcx
LOAD:00000000004F22F    retn
LOAD:00000000004F22F sub_44F1C5    endp ; sp-analysis Failed

```

0x44f214的rep指令把0x400000到0x44f248复制到0x800000位置。

0x44f22c处，rbp是0x84efb5。这句话是start处的第二条指令。这里不必步入。如果步入，

在44EFCB处的call sub\_44F020，再步过就会卡死（不知道为什么，奇怪！如果卡死，按ctrl+c可以只是终止而不退出gdb）；步入之后可以使用finish命令运行到函数外，即0x44F22E处pop rbp。

在0x44f22f retn之后，到达0x84f248: call 0x84f2ac。

0x84f2ac是一开始没有复制的地址，可以考虑是0x84efb5的代码形成的新代码。此处内存开辟是在系统调用mmap处，刚开辟后此处全为0。

此时可以dump出来0x800000处的内存，见dump2.（我修改了ep, va, filesize）

```

LOAD:000000000084F2D7
LOAD:000000000084F2D7 loc_84F2D7:           ; CODE XREF: start+33↓ j
LOAD:000000000084F2D7 cmp    qword ptr [rsi], 0
LOAD:000000000084F2D8 movsq
LOAD:000000000084F2D9 movsq
LOAD:000000000084F2D9 jnz    short loc_84F2D7
LOAD:000000000084F2E1 lea    r15, [rdi-8]
LOAD:000000000084F2E5 mov    [rdx], rdi
LOAD:000000000084F2E8 mov    eax, 3D202020h
LOAD:000000000084F2ED stosd
LOAD:000000000084F2EE mov    edx, 1000h      ; bufsiz
LOAD:000000000084F2F3 mov    rsi, rdi      ; buf
LOAD:000000000084F2F6 mov    rdi, r9       ; path
LOAD:000000000084F2F9 push   59h
LOAD:000000000084F2FB pop    rax
LOAD:000000000084F2FC syscall          ; LINUX - sys_readlink
LOAD:000000000084F2FE test   eax, eax
LOAD:000000000084F300 js    short loc_84F306
LOAD:000000000084F302 mov    byte ptr [rsi+rax], 0
LOAD:000000000084F306
LOAD:000000000084F306 loc_84F306:           ; CODE XREF: start+54↑ j
LOAD:000000000084F306 add    r9, 0Fh
LOAD:000000000084F30A pop    rcx
LOAD:000000000084F30B pop    rsi
LOAD:000000000084F30C pop    rdi
LOAD:000000000084F30D sub    rsp, 800h
LOAD:000000000084F314 mov    rdx, rsp
LOAD:000000000084F317 mov    r8, rbp
LOAD:000000000084F31A push   0
LOAD:000000000084F31C call   sub_84F782      ; 可以步过
LOAD:000000000084F321 pop    rdx
LOAD:000000000084F322 add    rsp, 800h
LOAD:000000000084F329 pop    rsi
LOAD:000000000084F32A pop    rdi
LOAD:000000000084F32B pop    rcx
LOAD:000000000084F32C pop    rcx
LOAD:000000000084F32D shr    ecx, 0Ch
LOAD:000000000084F330 add    rdi, rcx
LOAD:000000000084F333 sub    esi, ecx
LOAD:000000000084F335 push   rax
LOAD:000000000084F336 push   08h
LOAD:000000000084F338 pop    rax
LOAD:000000000084F339 jmp    qword ptr [r15]
LOAD:000000000084F339 start
LOAD:000000000084F339 endp

```

接着运行到0x84f31c: call 0x84f782, 可以步过。

在0x84f339: jmp QWORD PTR [r15], 跳转到0x40000c

此时可以dump处0x400000处的内存，见文件bindump.so。这就是脱壳后的文件了，但是不能运行，不清楚原因。可以看到ep是0x400890。

0x40000c: syscall (调用号11, sys\_munmap, 解除内存映射)

0x40000e: ret

上面两处是elf头部的padding中存储的，感叹做的壳之精准。

ret之后到达：

0x400890: xor ebp,ebp

下一部分，分析主逻辑。

Direction	Type	Address	Text
Up	o	sub_4009EF+4F	mov    esi, offset aEnterTheFlag; "Enter the flag:\n"
Down	o	LOAD:000000000004C8EC2	mov    rsi, offset aEnterTheFlag; "Enter the flag:\n"

OK Cancel Search Help

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发现Enter the flag出现了两次。

按照程序执行过程只会执行上面一个。

```
; Attributes: bp-based Frame

sub_4009EF proc near

var_78= qword ptr -78h
var_70= byte ptr -70h
var_8= qword ptr -8

push    rbp
mov     rbp, rsp
add    rsp, 0FFFFFFFFFFFF80h
; 10: v6 = __readfsqword(0x28u);
mov     rax, fs:28h
mov     [rbp+var_8], rax
; 11: if ( anti_debug_43F380(0LL, 0LL, 0LL, 0LL) )
xor    eax, eax
mov     ecx, 0
mov     edx, 0
mov     esi, 0
mov     edi, 0          ; request
mov     eax, 0
call    anti_debug_43F380
mov     [rbp+var_78], rax
cmp    [rbp+var_78], 0
jz     short loc_400A39
```

```
; 12:     sub_40EAD0(0LL);
mov     edi, 0
call    sub_40EAD0

; 13:     write_43E9B0(1LL, (__int64)"Enter the flag:\n");

loc_400A39:
mov     edx, 10h
mov     esi, offset aEnterTheFlag ; "Enter the flag:\n"
mov     edi, 1
call    write_43E9B0
; 14:     read_43E950(0LL, &v5, 32LL);
lea    rax, [rbp+var_70]
```

```

1 int64 sub_4009EF()
2 {
3     const char *v0; // rsi
4     _int64 v1; // rdx
5     _int64 result; // rax
6     _int64 v3; // rcx
7     unsigned __int64 v4; // rt1
8     char v5; // [rsp+10h] [rbp-70h]
9     unsigned __int64 v6; // [rsp+78h] [rbp-8h]
10
11    v6 = __readfsqword(0x28u);
12    if ( anti_debug_43F380(0LL, 0LL, 0LL, 0LL) )
13        sub_40EAD0(0LL);
14    write_43E9B0(1LL, (_int64)"Enter the flag:\n");
15    read_43E950(0LL, &v5, 32LL);
16    if ( (unsigned int)sub_4009AE((__int64)&v5) != 0 )
17    {
18        v0 = "You are right\n";
19        write_43E9B0(1LL, (_int64)"You are right\n");
20    }
21    else
22    {
23        v0 = "You are wrong\n";
24        write_43E9B0(1LL, (_int64)"You are wrong\n");
25    }
26    result = 0LL;
27    v4 = __readfsqword(0x28u);
28    v3 = v4 ^ v6;
29    if ( v4 != v6 )
30        sub_442480(1LL, v0, v1, v3);
31    return result;
32}

```

从流程图上可以看到左枝非正常退出，猜测和反调试有关，进入43f380函数，看到调用了ptrace的系统调用，确实是反调试。由于没有脱壳成功，这里没办法修改判断条件，只好用gdb脚本修改eflags寄存器了。

接着分析0x4009AE函数：

```

1 BOOL8 __fastcall sub_4009AE(__int64 a1)
2 {
3     return (unsigned int)sub_400360(a1, (_int64)"qwb{this_is_wrong_flag}") == 0;
4 }

```

这个时候就应该认识到这是错误的分支了，毕竟告诉你是错的flag了。但是我继续分析了，

```

1 int64 __fastcall sub_400360(__int64 a1, __int64 a2)
2 {
3     return qword_6C9850(a1, a2);
4 }

```

具体函数还没有写入。根据动态调试，最后比较字符串的函数是42D820.

```

33    __asm
34    {
35        movdqu xmm1, xmmword ptr [rdi]; 简单说就是mov
36        movdqu xmm0, xmmword ptr [rsi]
37        pcmpeqb xmm0, xmm1; 比较, 相同byte置ff. 字符串的首字节存在最后byte
38        pminub xmm0, xmm1; 感觉是逐字节取小的
39        pxor  xmm1, xmm1
40        pcmpeqb xmm0, xmm1
41        pmovmskb eax, xmm0; 取每个字节的最高位
42    }
43    if (_RAX)
44        goto LABEL_19;
45    __asm
46    {
47        movdqu xmm6, xmmword ptr [rdi+10h]; 自己输入
48        movdqu xmm3, xmmword ptr [rsi+10h]
49        movdqu xmm5, xmmword ptr [rdi+20h]
50        pcmpeqb xmm3, xmm6
51        movdqu xmm2, xmmword ptr [rsi+20h]
52        pminub xmm3, xmm6
53        pcmpeqb xmm3, xmm1; xmm1 是0
54        movdqu xmm4, xmmword ptr [rdi+30h]
55        pcmpeqb xmm2, xmm5
56        pmovmskb edx, xmm3
57        movdqu xmm0, xmmword ptr [rsi+30h]
58        pminub xmm2, xmm5
59        pcmpeqb xmm2, xmm1
60        pcmpeqb xmm0, xmm4
61        pmovmskb eax, xmm2
62        pminub xmm0, xmm4
63        pcmpeqb xmm0, xmm1
64        pmovmskb ecx, xmm0
65    }

```

仔细分析逻辑，需要保证输入地址的64字节和目标地址的64字节相等，而只能输入32字节，根本不能实现。

废了很久时间之后，想到了另一处Enter the flag。

```

LOAD:00000000004C8E92          align 20h
LOAD:00000000004C8EA0          xor    rdi, rdi
LOAD:00000000004C8EA3          xor    rsi, rsi
LOAD:00000000004C8EA6          xor    rdx, rdx
LOAD:00000000004C8EA9          xor    r10, r10
LOAD:00000000004C8EAC          mov    eax, 65h
LOAD:00000000004C8EB1          syscall           ; LINUX - sys_trace
LOAD:00000000004C8EB3          cmp    eax, 0
LOAD:00000000004C8EB6          jnz    locret_4C8FDB
LOAD:00000000004C8EBC          xor    rdi, rdi
LOAD:00000000004C8EBF          inc    rdi
LOAD:00000000004C8EC2          mov    rsi, offset aEnterTheFlag ; "Enter the flag:\n"
LOAD:00000000004C8EC9          mov    rdx, 10h
LOAD:00000000004C8ED0          xor    eax, eax
LOAD:00000000004C8ED2          inc    eax
LOAD:00000000004C8ED4          syscall           ; LINUX - sys_write
LOAD:00000000004C8ED6          xor    rdi, rdi
LOAD:00000000004C8ED9          xor    eax, eax
LOAD:00000000004C8EDB          mov    rsi, offset input_6CCDB0
LOAD:00000000004C8EE2          mov    rdx, 20h
LOAD:00000000004C8EE9          syscall           ; LINUX - sys_read
LOAD:00000000004C8EEB          cmp    eax, 0
LOAD:00000000004C8EEE          jle    loc_4C8FA9
LOAD:00000000004C8EF4 ; 7:   if ( strlen(input_6CCDB0) == 21
LOAD:00000000004C8EF4 ; 8:     && input_6CCDB0[1] == 'w'
LOAD:00000000004C8EF4 ; 9:     && input_6CCDB0[2] == 'b'
LOAD:00000000004C8EF4 ; 10:    && input_6CCDB0[3] == '{'
LOAD:00000000004C8EF4 ; 11:    && input_6CCDB0[20] == '}' )
LOAD:00000000004C8EF4
LOAD:00000000004C8EF4 ; ===== S U B R O U T I N E =====
LOAD:00000000004C8EF4
LOAD:00000000004C8EF4
LOAD:00000000004C8EF4 sub_4C8EF4 proc near
LOAD:00000000004C8EF4     mov    rdi, offset input_6CCDB0
LOAD:00000000004C8EF4     mov    rcx, 0xFFFFFFFFFFFFFFh
LOAD:00000000004C8FB     xor    eax, eax
LOAD:00000000004C8F02    repne scasb
LOAD:00000000004C8F04    not    rcx
LOAD:00000000004C8F06    sub    rcx, 1
LOAD:00000000004C8F09

```

这里并未被ida识别成函数，0x4c8ef4处是我标记的函数，从这里标记可以使函数能使用f5.由于原程序并不会执行到此处，需要gdb脚本中修改pc值。

```
1 signed __int64 sub_4C8EF4()
2 {
3     _BYTE *v0; // rdi
4     __int64 *v1; // rsi
5     unsigned __int64 v2; // rdx
6     signed __int64 result; // rax
7
8     if ( strlen(input_6CCDB0) == 21
9         && input_6CCDB0[1] == 'w'
10        && input_6CCDB0[2] == 'b'
11        && input_6CCDB0[3] == '{'
12        && input_6CCDB0[20] == '}' )
13    {
14         encrypt1_4C8CC0((__int64)&input_4_6CCDB4);
15         sub_4C8E50((__int64)&input_4_6CCDB4);
16         encrypt1_4C8CC0((__int64)&input_4_6CCDB4);
17         sub_4C8E50((__int64)&input_4_6CCDB4);
18         encrypt1_4C8CC0((__int64)&input_4_6CCDB4);
19         v0 = &input_4_6CCDB4;
20         sub_4C8E50((__int64)&input_4_6CCDB4);
21         v1 = qword_4C8CB0; // target
22         v2 = 0LL;
23         while ( v2 < 0x10 && *v0 == *(BYTE *)v1 )
24     {
25             ++v2;
26             ++v0;
27             v1 = (__int64 *)((char *)v1 + 1);
28     }
29 }
30     __asm { syscall; LINUX - sys_write }
31     result = 60LL;
32     __asm { syscall; LINUX - sys_exit }
33     return result;
34 }
```

下面是可以快速进入关键位置的gdb脚本。

```

#!/bin/bash
file hide
b *0x44f22f
r
si
si
d
b *0x84f339
#jmp    QWORD PTR [r15];0x40000c
c
d
si
si
si
b* 0x4009EF
c
set $rip=0x4C8EBC
b *0x4c8f11
c
set $eflags = $eflags |(1<<6)
5
b *0x4C8D09
c
c
d
b *0x4C8DFB
c

```

下面是本题目的exp

```

#coding=utf-8
import struct
import string
def u32(data):
    return struct.unpack("<I",data)[0]

def p32(data):
    return struct.pack("<I",data)

def u64(data):
    return struct.unpack("<Q",data)[0]

def p64(data):
    return struct.pack("<Q",data)

input1 = '1234567890123456'
input1 = bytearray(input1)
CONST_STR = 's1IpP3rEv3Ryd4Y3'
CONST = 0X676E696C
v4 = 0
v4_4=0
v4_4_arr = [0 for i in range(0,9)]
for i in range(1,9):
    v4_4_arr[i] = (v4_4_arr[i-1]+CONST)&0xFFFFFFFF

def re_block(byte_arr_8):
    . .

```

```

i = 0
v3 = u32(byte_arr_8[0:4])
v4 = u32(byte_arr_8[4:8])
print 'round', v3, v4
for i in range(7,-1,-1):
    v30 = (v3 << 4) & 0xffffffff
    v2c = v3 >> 5
    edx = v30 ^ v2c
    v30 = (v3 + edx) & 0xffffffff

    v28 = (v4_4_arr[i+1] >> 11) & 3
    edx = u32(CONST_STR[v28 * 4:(v28 + 1) * 4])
    v2c = (v4_4_arr[i+1] + edx) & 0xffffffff # xxxx
    # print 'v2c',hex(v2c)
    print v30 ^ v2c
    v4 = (v4+0x100000000-(v30 ^ v2c)) & 0xffffffff

    v30 = (v4 << 4) & 0xffffffff
    v2c = v4 >> 5
    edx = v30 ^ v2c

    v30 = (v4 + edx) & 0xffffffff

    v28 = v4_4_arr[i] & 3
    edx = u32(CONST_STR[v28 * 4:(v28 + 1) * 4])
    v2c = (v4_4_arr[i] + edx) & 0xffffffff

    v3 = (v3+0x100000000-(v30 ^ v2c)) & 0xffffffff

    print 'round',i,v3,v4
byte_arr_8[0:4] = p32(v3)
byte_arr_8[4:8] = p32(v4)
return byte_arr_8

def xor16(byte_arr_16):
    for i in range(0,16):
        byte_arr_16[i]^=i

def re_all(str16):
    byte_arr = bytarray(str16)
    xor16(byte_arr)#传入整个bytarray, 就是传入地址
    byte_arr[0:8] = re_block(byte_arr[0:8])#传入部分bytarray, 就是复制之后再传入
    byte_arr[8:16] = re_block(byte_arr[8:16])
    xor16(byte_arr)
    byte_arr[0:8] = re_block(byte_arr[0:8])
    byte_arr[8:16] = re_block(byte_arr[8:16])
    xor16(byte_arr)
    byte_arr[0:8] = re_block(byte_arr[0:8])
    byte_arr[8:16] = re_block(byte_arr[8:16])
    return str(byte_arr)

def block(str8):
    i=0
    input1=bytarray(str8)
    v3 = u32(input1[8 * i:8 * i + 4])
    v4 = u32(input1[8 * i + 4:8 * (i + 1)])
    v4_4 = 0 ##0000
    for j in range(0, 8):

        v30 = (v4 << 4) & 0xffffffff
        v2c = v4 >> 5

```

```

edx = v30 ^ v2c

v30 = (v4 + edx) & 0xffffffff

v28 = v4_4 & 3
edx = u32(CONST_STR[v28 * 4:(v28 + 1) * 4])
v2c = (v4_4 + edx) & 0xffffffff

v3 = ((v30 ^ v2c) + v3) & 0xffffffff

v4_4 = (v4_4 + CONST) & 0xffffffff

v30 = (v3 << 4) & 0xffffffff
v2c = v3 >> 5
edx = v30 ^ v2c
v30 = (v3 + edx) & 0xffffffff

v28 = (v4_4 >> 11) & 3
edx = u32(CONST_STR[v28 * 4:(v28 + 1) * 4])
v2c = (v4_4 + edx) & 0xffffffff # xxxx
print v30 ^ v2c
v4 = ((v30 ^ v2c) + v4) & 0xffffffff

print 'round', j, v3, v4

input1[8 * i:8 * i + 4] = p32(v3)
input1[8 * i + 4:8 * (i + 1)] = p32(v4)
str8_1=str(input1)
return str8_1

def block2(str8):
    input1 = bytearray(str8)
    for i in range(0,8):
        input1[i]=input1[1]^i
    return str(input1)

des = ('52B8137F358CF21B'+ 'F46386D2734F1E31').decode('hex')
print len(des)

print block('12345678').encode('hex')
des1 = '5b90ef3f91b58fe6'.decode('hex')
print re_all(bytearray(des))

```